

CYCLIC CODES OVER THE RING

$\mathbb{F}_p[u, v, w]/\langle u^2, v^2, w^2, uv - vu, vw - wv, uw - wu \rangle$

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ABSTRACT. Let $R_{u^2, v^2, w^2, p}$ be a finite non chain ring $\mathbb{F}_p[u, v, w]/\langle u^2, v^2, w^2, uv - vu, vw - wv, uw - wu \rangle$, where p is a prime number. This ring is a part of family of Frobenius rings. In this paper, we explore the structures of cyclic codes over the ring $R_{u^2, v^2, w^2, p}$ of arbitrary length. We obtain a unique set of generators for these codes and also characterize free cyclic codes. We show that Gray images of cyclic codes are 8-quasicyclic binary linear codes of length $8n$ over \mathbb{F}_p . We also determine the rank and the Hamming distance for these codes. At last, we have given some examples.

1. Introduction

Cyclic codes are a key family of linear codes because of their lavish algebraic structures and practical accomplishment. A considerable attention has been paid to cyclic codes over rings in the early 1990's because of their affluent applications to design error-correcting coding schemes for wireless communication system. The classification of cyclic codes of length n over chain rings as well as over some non chain rings, when $\gcd(n, p) = 1$, where p is the characteristic of the ring, has been completely discussed in the literature. On the other hand, the classification of codes when p divides n is still not complete. Cyclic codes have been extensively studied over various finite chain rings in [1–5, 10]. More latterly, cyclic codes over finite non-chain rings have also been contemplated. However, the analysis on non-chain rings seems to be challenging as the algebraic structure does not allow to give nice and compact presentation of linear codes over these rings. Yildiz and Karadeniz in [12] characterized cyclic codes of odd length over the non-chain ring $\mathbb{F}_2[u, v]/\langle u^2, v^2, uv - vu \rangle$ and obtained some good binary codes under two Gray maps. Sobhani and Molakarimi in [11] characterized the cyclic codes over the ring $F_{2^m}[u, v]/\langle u^2, v^2, uv - vu \rangle$, they also determined \mathbb{F}_{2^m} -basis as well as the mass the formula for the number of these codes. The authors of [8] extended these studies in more general way for cyclic

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codes over the ring $\mathbb{F}_p[u, v]/\langle u^2, v^2, uv - vu \rangle$ and presented some ternary optimal codes as gray images. In [6], S. T. Dougherty et al. considered a family of rings $F_2[u_1, u_2, \dots, u_k]/\langle u_1^2, u_2^2, \dots, u_k^2, u_i u_j - u_j u_i \rangle$ and studied one-generator cyclic codes.

The purpose of this paper is to obtain the structure theorems for cyclic codes over the non-chain $R_{u^2, v^2, w^2, p}$ ring in more general setting. The paper is organized as follows: In Section 2, we find a generating set of polynomials along with the conditions on these generators for the cyclic codes over the ring $R_{u^2, v^2, w^2, p}$ of length n . We also discuss here generating polynomials for cases of free cyclic codes and n relatively prime to p . In Section 3, we prove some lemmas, in which we express certain types of polynomials over $R_{u^2, v^2, w^2, p}$ as linear combinations of generators of cyclic codes over $R_{u^2, v^2, w^2, p}$. By help of these lemmas, we derive a minimal spanning set, we also calculate ranks of these codes. In Section 4, we obtain the minimum distance of corresponding codes of length p^l . In Section 5.1, we will show that gray images of cyclic codes are 8-quasicyclic binary linear code of length $8n$ over \mathbb{F}_p . In Section 6, we discuss some examples of cyclic codes over the ring $R_{u^2, v^2, w^2, p}$ of length 3, 4, 5 over $\mathbb{F}_3, \mathbb{F}_2, \mathbb{F}_5$ respectively.

2. The structures of cyclic codes over the ring $R_{u^2, v^2, w^2, p}$

Let $R_{u^2, v^2, w^2, p} = \mathbb{F}_p[u, v, w]/\langle u^2, v^2, w^2, uv - vu, vw - wv, uw - wu \rangle$, where p is a prime number and n is a positive integer. We can write $R_{u^2, v^2, w^2, p}$ as $R_{u^2, v^2, w^2, p} = R_{u^2, v^2, p} + wR_{u^2, v^2, p}$, $w^2 = 0$, where $R_{u^2, v^2, p} = \mathbb{F}_p + u\mathbb{F}_p + v\mathbb{F}_p + uv\mathbb{F}_p$ and $u^2 = 0, v^2 = 0$. Let $R_{u^2, v^2, w^2, p, n} = R_{u^2, v^2, w^2, p}[x]/\langle x^n - 1 \rangle$. Let C be a cyclic code of length n over the ring $R_{u^2, v^2, w^2, p}$. We can also consider C as an ideal in the ring $R_{u^2, v^2, w^2, p, n}$. We define the map $\psi : R_{u^2, v^2, w^2, p} \rightarrow R_{u^2, v^2, p}$ by $\psi(\alpha + w\beta) = \alpha$, where $\alpha, \beta \in R_{u^2, v^2, p}$. Clearly the map ψ is a surjective ring homomorphism. Let $R_{u^2, v^2, p, n} = R_{u^2, v^2, p}[x]/\langle x^n - 1 \rangle$. We extend this homomorphism to a homomorphism ϕ from C to the ring $R_{u^2, v^2, p, n}$ defined by

$$(2.1) \quad \phi(c_0 + c_1x + \dots + c_{n-1}x^{n-1}) = \psi(c_0) + \psi(c_1)x + \dots + \psi(c_{n-1})x^{n-1},$$

where $c_i \in R_{u^2, v^2, w^2, p}$. Let $J = \{r(x) \in R_{u^2, v^2, p, n}[x] : wr(x) \in \ker\phi\}$. We see that J is an ideal of $R_{u^2, v^2, p, n}$. Hence, we can consider J as a cyclic code over $R_{u^2, v^2, p}$. We know from Theorem 3.1 of [8] that any ideal of $R_{u^2, v^2, p, n}$ is of the form $\langle g(x) + up_1(x) + vq_1(x) + uvr_1(x), ua_1(x) + vq_2(x) + uvr_2, va_2(x) + uvr_3(x), wa_3(x) \rangle$. Now we assume that $B_1 = g(x) + up_1(x) + vq_1(x) + uvr_1(x)$, $B_2 = ua_1(x) + vq_2(x) + uvr_2$, $B_3 = va_2(x) + uvr_3(x)$, $B_4 = wa_3(x)$. So $J = \langle B_1, B_2, B_3, B_4 \rangle$. Therefore, we can write $\ker\phi = \langle wB_1, wB_2, wB_3, wB_4 \rangle$. Since ϕ is a surjective homomorphism, the image $\text{Im}\phi$ is an ideal of $R_{u^2, v^2, p, n}$. Hence, $\text{Im}\phi$ is a cyclic code over $R_{u^2, v^2, p}$. Again we can write $\text{Im}\phi$ as above. That is, $\text{Im}\phi = \langle B'_1, B'_2, B'_3, B'_4 \rangle$. Therefore, the code C over the ring $R_{u^2, v^2, w^2, p}$ can be written as $C = \langle A_1, A_2, \dots, A_8 \rangle$, where, A_i 's are defined as follows:

$$A_1 = f_1(x) + uf_{1,2}(x) + vf_{1,3}(x) + uvf_{1,4}(x) + wf_{1,5}(x) + uwf_{1,6}(x)$$

$$\begin{aligned}
& + vwf_{1,7}(x) + uvwf_{1,8}(x), \\
A_2 &= uf_2(x) + vf_{2,3}(x) + uvf_{2,4}(x) + wf_{2,5}(x) + uvwf_{2,6}(x) + vwf_{2,7}(x) \\
& + uvwf_{2,8}(x), \\
A_3 &= vf_3(x) + uvf_{3,4}(x) + wf_{3,5}(x) + uvwf_{3,6}(x) + vwf_{3,7}(x) + uvwf_{3,8}(x), \\
A_4 &= uvf_4(x) + wf_{4,5}(x) + uvwf_{4,6}(x) + vwf_{4,7}(x) + uvwf_{4,8}(x), \\
A_5 &= wf_5(x) + uvwf_{5,6}(x) + vwf_{5,7}(x) + uvwf_{5,8}(x), \\
A_6 &= uvwf_6(x) + vwf_{6,7}(x) + uvwf_{6,8}(x), \\
A_7 &= vwf_7(x) + uvwf_{7,8}(x), \\
A_8 &= uvwf_8(x).
\end{aligned}$$

Throughout this paper we use A_1, A_2, \dots, A_8 for above polynomials.

For an ideal C of the ring $R_{u^2, v^2, w^2, p, n} = R_{u^2, v^2, w^2, p}[x]/\langle x^n - 1 \rangle$, we define the residue and the torsion of an ideal C as

$$\begin{aligned}
\text{Res}(C) &= \{a \in R_{u^2, v^2, p, n} \mid \exists b \in R_{u^2, v^2, p, n} : a + wb \in C\} \text{ and} \\
\text{Tor}(C) &= \{a \in R_{u^2, v^2, p, n} \mid wa \in C\}.
\end{aligned}$$

It can be easily shown that when C is an ideal of $R_{u^2, v^2, w^2, p, n}$, $\text{Res}(C)$ and $\text{Tor}(C)$ both are ideals of $R_{u^2, v^2, p, n}$. And it is easy to show that $\text{Res}(C) = \text{Im}\phi$ and $\text{Tor}(C) = J$. Now we define eight ideals associated to C .

$$\begin{aligned}
(2.2) \quad C_1 &= \text{Res}(\text{Res}(\text{Res}(C))) \\
&= C \text{ mod } \langle u, v, w \rangle = \langle f_1(x) \rangle, \\
(2.3) \quad C_2 &= \text{Tor}(\text{Res}(\text{Res}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid uf(x) \in C \text{ mod } \langle v, w \rangle\} = \langle f_2(x) \rangle, \\
(2.4) \quad C_3 &= \text{Res}(\text{Tor}(\text{Res}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid vf(x) \in C \text{ mod } \langle uv, w \rangle\} = \langle f_3(x) \rangle, \\
(2.5) \quad C_4 &= \text{Tor}(\text{Tor}(\text{Res}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid uvf(x) \in C \text{ mod } \langle w \rangle\} = \langle f_4(x) \rangle, \\
(2.6) \quad C_5 &= \text{Res}(\text{Res}(\text{Tor}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid wf(x) \in C \text{ mod } \langle uw, vw \rangle\} = \langle f_5(x) \rangle, \\
(2.7) \quad C_6 &= \text{Tor}(\text{Res}(\text{Tor}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid uvf(x) \in C \text{ mod } \langle vw \rangle\} = \langle f_6(x) \rangle, \\
(2.8) \quad C_7 &= \text{Res}(\text{Tor}(\text{Tor}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid vwf(x) \in C \text{ mod } \langle uvw \rangle\} = \langle f_7(x) \rangle, \\
(2.9) \quad C_8 &= \text{Tor}(\text{Tor}(\text{Tor}(C))) \\
&= \{f(x) \in \mathbb{F}_p[x] \mid uvwf(x) \in C\} = \langle f_8(x) \rangle.
\end{aligned}$$

These are ideals of $\mathbb{F}_p[x]/\langle x^n - 1 \rangle$, hence principal ideals. Throughout this paper we use C_1, C_2, \dots, C_8 for above ideals.

Theorem 1. *Any ideal C of the ring $R_{u^2, v^2, w^2, p, n}$ is uniquely generated by the polynomials A_1, A_2, \dots, A_8 with $f_{i,j}(x) = 0$ or $\deg(f_{i,j}(x)) < \deg(f_j(x))$, where A_i, f_i and $f_{i,j}$ are defined as above.*

Proof. Proof is similar to the proof of Theorem 1 [11]. \square

Theorem 2. *Let $C = \langle A_1, A_2, \dots, A_8 \rangle$ be an ideal of the ring $R_{u^2, v^2, w^2, p, n}$. Then we must have*

- (1) $f_8(x)|f_i(x)$ for $1 \leq i \leq 7$; $f_j(x)|f_1(x)|(x^n - 1)$ for $2 \leq j \leq 7$;
- (2) $f_4(x)|f_2(x)$; $f_4(x)|f_3(x)$; $f_6(x)|f_5(x)$; $f_6(x)|f_2(x)$; $f_7(x)|f_5(x)$;
 $f_7(x)|f_3(x)$;
- (3) $f_{i+1}(x)|f_{i+1}(x)\left(\frac{x^n-1}{f_i(x)}\right)$ for $1 \leq i \leq 7$;
- (4) For a fix j , $1 \leq j \leq 7$, $f_{i+j}(x)|\frac{x^n-1}{f_i(x)}\frac{x^n-1}{f_{i+1}(x)}\dots\frac{x^n-1}{f_{i+j-1}(x)}f_{i,i+j}(x)$ for
 $1 \leq i \leq 8 - j$;
- (5) $f_i(x)|\frac{x^n-1}{f_{i-2}(x)}\left(f_{i-2,i}(x) - \frac{f_{i-2,i-1}(x)}{f_{i-1}(x)}f_{i-1,i}(x)\right)$ for $3 \leq i \leq 8$;
- (6) $f_i(x)|\frac{x^n-1}{f_{i-3}(x)}\left(f_{i-3,i}(x) - \frac{f_{i-3,i-2}(x)}{f_{i-2}(x)}f_{i-2,i}(x)\right.$
 $\left. - \frac{f_{i-3,i-1}(x) - \frac{f_{i-3,i-2}(x)}{f_{i-2}(x)}f_{i-2,i-1}(x)}{f_{i-1}(x)}f_{i-1,i}(x)\right)$ for $4 \leq i \leq 8$;
- (7) $f_i(x)|\frac{x^n-1}{f_{i-4}(x)}\left(f_{i-4,i}(x) - \frac{f_{i-4,i-3}(x)}{f_{i-3}(x)}f_{i-3,i}(x) - Af_{i-2,i}(x) - Bf_{i-1,i}(x)\right)$
for $i \in (5, 6, 7, 8)$, where
 $A = \left(\frac{f_{i-4,i-2}(x) - \frac{f_{i-4,i-3}(x)}{f_{i-3}(x)}f_{i-3,i-2}(x)}{f_{i-2}(x)}\right)$ and
 $B = \left(\frac{f_{i-4,i-1}(x) - \frac{f_{i-4,i-3}(x)}{f_{i-3}(x)}f_{i-3,i-1}(x) - Af_{i-2,i-1}(x)}{f_{i-1}(x)}\right)$.
- (8) $f_i(x)|\frac{x^n-1}{f_{i-5}(x)}\left(f_{i-5,i}(x) - \frac{f_{i-5,i-4}(x)}{f_{i-4}(x)}f_{i-4,i}(x) - Af_{i-3,i}(x)\right.$
 $\left. - Bf_{i-2,i}(x) - Df_{i-1,i}(x)\right)$ for $i \in (6, 7, 8)$, where
 $A = \left(\frac{f_{i-5,i-3}(x) - \frac{f_{i-5,i-4}(x)}{f_{i-4}(x)}f_{i-4,i-3}(x)}{f_{i-3}(x)}\right)$,
 $B = \left(\frac{f_{i-5,i-2}(x) - \frac{f_{i-5,i-4}(x)}{f_{i-4}(x)}f_{i-4,i-2}(x) - Af_{i-3,i-2}(x)}{f_{i-2}(x)}\right)$ and
 $D = \left(\frac{f_{i-5,i-1}(x) - \frac{f_{i-5,i-4}(x)}{f_{i-4}(x)}f_{i-4,i-1}(x) - Af_{i-3,i-1}(x) - Bf_{i-2,i-1}(x)}{f_{i-1}(x)}\right)$.
- (9) $f_i(x)|\frac{x^n-1}{f_{i-6}(x)}\left(f_{i-6,i}(x) - Af_{i-5,i}(x) - Bf_{i-4,i}(x) - Df_{i-3,i}(x)\right.$
 $\left. - Ef_{i-2,i}(x) - Ff_{i-1,i}(x)\right)$ for $i \in (7, 8)$, where
 $A = \left(\frac{f_{i-6,i-5}(x)}{f_{i-5}(x)}\right)$,
 $B = \left(\frac{f_{i-6,i-4}(x) - Af_{i-5,i-4}(x)}{f_{i-4}(x)}\right)$,

$$\begin{aligned}
D &= \left(\frac{f_{i-6,i-3}(x) - Af_{i-5,i-3}(x) - Bf_{i-4,i-3}(x)}{f_{i-3}(x)} \right), \\
E &= \left(\frac{f_{i-6,i-2}(x) - Af_{i-5,i-2}(x) - Bf_{i-4,i-2}(x) - Df_{i-3,i-2}(x)}{f_{i-2}(x)} \right) \text{ and} \\
F &= \left(\frac{f_{i-6,i-1}(x) - Af_{i-5,i-1}(x) - Bf_{i-4,i-1}(x) - Df_{i-3,i-1}(x) - Ef_{i-2,i-1}(x)}{f_{i-1}(x)} \right). \\
(10) \quad f_8(x) &| \frac{x^n-1}{f_1(x)} \left(f_{1,8}(x) - Af_{2,8}(x) - Bf_{3,8}(x) - Df_{4,8}(x) - Ef_{5,8}(x) \right. \\
&\quad \left. - Ff_{6,8}(x) - Gf_{7,8}(x) \right), \text{ where} \\
A &= \left(\frac{f_{1,2}(x)}{f_2(x)} \right), \\
B &= \left(f_{1,3}(x) - Af_{2,3}(x) \right), \\
D &= \left(\frac{f_{1,4}(x) - Af_{2,4}(x) - Bf_{3,4}(x)}{f_4(x)} \right), \\
E &= \left(\frac{f_{1,5}(x) - Af_{2,5}(x) - Bf_{3,5}(x) - Df_{4,5}(x)}{f_5(x)} \right), \\
F &= \left(\frac{f_{1,6}(x) - Af_{2,6}(x) - Bf_{3,6}(x) - Df_{4,6}(x) - Ef_{5,6}(x)}{f_6(x)} \right) \text{ and} \\
G &= \left(\frac{f_{1,7}(x) - Af_{2,7}(x) - Bf_{3,7}(x) - Df_{4,7}(x) - Ef_{5,7}(x) - Ff_{6,7}(x)}{f_7(x)} \right). \\
(11) \quad f_i(x) &| f_{i-2,i-1}(x) \text{ for } i \in (4, 6, 8); \\
(12) \quad f_i(x) &| \left(f_{1,2}(x) - \frac{f_1(x)}{f_{i-1}(x)} f_{i-1,i}(x) \right) \text{ for } i \in (4, 6, 8); \\
(13) \quad f_i(x) &| \left(f_{i-5,i-4}(x) - \frac{f_{i-5}(x)}{f_{i-1}(x)} f_{i-1,i}(x) \right) \text{ for } i \in (7, 8); \\
(14) \quad f_i(x) &| \left(f_{i-6,i-4}(x) - \frac{f_{i-6}(x)}{f_{i-2}(x)} f_{i-2,i}(x) \right. \\
&\quad \left. + \frac{f_{i-6,i-5}(x) - \frac{f_{i-6}(x)}{f_{i-2}(x)} f_{i-2,i-1}(x)}{f_{i-1}(x)} f_{i-1,i}(x) \right) \text{ for } i \in (7, 8); \\
(15) \quad f_7(x) &| f_{4,5}(x) \text{ and } f_7(x) | f_{3,5}(x); \\
(16) \quad f_8(x) &| f_{2,5}(x); \\
(17) \quad f_8(x) &| \left(f_{3,6}(x) - \frac{f_{3,5}(x)}{f_7(x)} f_{7,8}(x) \right); \\
(18) \quad f_8(x) &| \left(f_{4,6}(x) - \frac{f_{4,5}(x)}{f_7(x)} f_{7,8}(x) \right); \\
(19) \quad f_8(x) &| \left(f_{5,6}(x) - \frac{f_5(x)}{f_7(x)} f_{7,8}(x) \right); \\
(20) \quad f_8(x) &| \left(f_{1,4}(x) - \frac{f_1(x)}{f_5(x)} f_{5,8}(x) - Af_{6,8}(x) - Bf_{7,8}(x) \right), \text{ where} \\
A &= \left(\frac{f_{1,2}(x) - \frac{f_1(x)}{f_5(x)} f_{5,6}(x)}{f_6(x)} \right) \text{ and } B = \left(\frac{f_{1,3}(x) - \frac{f_1(x)}{f_5(x)} f_{5,7}(x) - (A)f_{6,7}(x)}{f_6(x)} \right).
\end{aligned}$$

Proof. (1) We have $vwA_2 \in C$. Therefore, $uvwf_2(x) \in C$. This gives $f_2(x) \in C_8 = \langle f_8(x) \rangle$. Thus, $f_8(x) | f_2(x)$. Similarly, if we take uwA_3, wA_4, uvA_5, vA_6 and uA_7 we get $f_8(x) | f_i(x)$ for $3 \leq i \leq 7$.

(2) We have $vA_2 \in C$. Therefore, $uvf_2(x) \in C \text{ mod } w$. This gives $f_2(x) \in C_4 = \langle f_4(x) \rangle$. Thus, $f_4(x) | f_2(x)$. Similarly, if we take $uA_3, uA_5, wA_2, vA_5, wA_3$ and take mod by w, vw, vw, uvw, uvw respectively, we get the other conditions of (2).

(3) For $1 \leq i \leq 7$, we have $\frac{x^n-1}{f_i(x)}A_i \in C$. Therefore, $\frac{x^n-1}{f_i(x)}f_{i,i+1}(x) \in C_{i+1} = \langle f_{i+1}(x) \rangle$. Hence, $f_{i+1}(x) | \frac{x^n-1}{f_i(x)}f_{i,i+1}(x)$.

(4) For $j = 1$, Condition 4 is reduced to Condition 3. For $j = 2$ and for $1 \leq i \leq 6$, we have $\frac{x^n-1}{f_i(x)}\frac{x^n-1}{f_{i+1}(x)}A_i \in C$. This with Condition 3 gives $\frac{x^n-1}{f_i(x)}\frac{x^n-1}{f_{i+1}(x)}f_{i,i+2} \in C_{i+2} = \langle f_{i+2}(x) \rangle$. Hence, $f_{i+2}(x) | \frac{x^n-1}{f_i(x)}\frac{x^n-1}{f_{i+1}(x)}f_{i,i+2}$. This proves the condition for $j = 2$. Similarly for other value of j we can prove Condition 4.

(5) For $i = 3$, we have

$$\begin{aligned} & \left(\frac{x^n-1}{f_1(x)}A_1 - \frac{x^n-1}{f_1(x)}\frac{f_{1,2}(x)}{f_2(x)}A_2 \right) \\ &= v \frac{x^n-1}{f_1(x)}(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)) + uv \frac{x^n-1}{f_1(x)}(f_{1,4}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,4}(x)) \\ & \quad + w \frac{x^n-1}{f_1(x)}(f_{1,5}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,5}(x)) + uvw \frac{x^n-1}{f_1(x)}(f_{1,6}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,6}(x)) \\ & \quad + vvw \frac{x^n-1}{f_1(x)}(f_{1,7}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,7}(x)) + uvvw \frac{x^n-1}{f_1(x)}(f_{1,8}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,8}(x)) \\ & \in C. \end{aligned}$$

Since $v \frac{x^n-1}{f_1(x)}(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)) \in C \pmod{\langle uv, w \rangle}$, therefore,

$$\frac{x^n-1}{f_1(x)}(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)) \in C_3 \Rightarrow f_3(x) | \left(\frac{x^n-1}{f_1(x)}(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)) \right).$$

Similarly we get the results for rest of the values of i .

(6) For $i = 4$, we have

$$\left(\frac{x^n-1}{f_1(x)}A_1 - \frac{x^n-1}{f_1(x)}\frac{f_{1,2}(x)}{f_2(x)}A_2 + \frac{x^n-1}{f_1(x)f_3(x)}\left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)\right)A_3 \right) \in C.$$

Since $uv \frac{x^n-1}{f_1(x)}\left(f_{1,4}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,4}(x) + \frac{f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)}{f_3(x)}f_{3,4}(x)\right) \in C \pmod{w}$, therefore,

$$\begin{aligned} & \frac{x^n-1}{f_1(x)}\left(f_{1,4}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,4}(x) + \frac{f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)}{f_3(x)}f_{3,4}(x)\right) \in C_4 \\ & \Rightarrow f_4(x) | \frac{x^n-1}{f_1(x)}\left(f_{1,4}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,4}(x) + \frac{f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)}{f_3(x)}f_{3,4}(x)\right). \end{aligned}$$

Similarly we get the results for rest of the values of i .

(7) For $i = 5$, we have

$$\begin{aligned} & \frac{x^n-1}{f_1(x)}\left(A_1 - f_{1,2}(x)\frac{A_2}{f_2(x)} + \left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)\right)\frac{A_3}{f_3(x)}\right) \\ & - \left(\frac{x^n-1}{f_1(x)}\left(f_{1,4}(x) - f_{1,2}(x)\frac{A_2}{f_2(x)} + \left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)\right)\frac{A_3}{f_3(x)}\right)\frac{A_4}{f_4(x)}\right) \in C. \end{aligned}$$

Since

$$\begin{aligned} & w \frac{x^n-1}{f_1(x)}\left(f_{1,5}(x) - f_{1,2}(x)\frac{f_{2,5}(x)}{f_2(x)} + \left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)\right)\frac{f_{3,5}(x)}{f_3(x)}\right) \\ & - \frac{x^n-1}{f_1(x)}\left(f_{1,4}(x) - f_{1,2}(x)\frac{f_{2,4}(x)}{f_2(x)} + \left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)}f_{2,3}(x)\right)\frac{f_{3,4}(x)}{f_3(x)}\right)\frac{f_{4,5}(x)}{f_4(x)} \end{aligned}$$

$\in C \pmod{(uw, vw)}$, therefore,

$$\begin{aligned} & \frac{x^n-1}{f_1(x)} \left(f_{1,5}(x) - f_{1,2}(x) \frac{f_{2,5}(x)}{f_2(x)} + \left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)} f_{2,3}(x) \right) \frac{f_{3,5}(x)}{f_3(x)} \right) \\ & - \frac{x^n-1}{f_1(x)} \left(f_{1,4}(x) - f_{1,2}(x) \frac{f_{2,4}(x)}{f_2(x)} + \left(f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)} f_{2,3}(x) \right) \frac{f_{3,4}(x)}{f_3(x)} \right) \frac{f_{4,5}(x)}{f_4(x)} \end{aligned}$$

$\in C_5$.

Similarly, we get the results for rest of the values of i .

$$(8) \text{ For } i = 6, \text{ we have } \left(\frac{x^n-1}{f_1(x)} A_1 - AA_2 - BA_3 + DA_4 + EA_5 \right) \in C$$

Since $uw \left(\frac{x^n-1}{f_1(x)} f_{1,6}(x) - Af_{2,6}(x) - Bf_{3,6}(x) + Df_{4,6}(x) + Ef_{5,6}(x) \right) \in C \pmod{vw}$, therefore,

$$\begin{aligned} & \frac{x^n-1}{f_1(x)} \left(f_{1,6}(x) - \frac{f_{1,2}(x)}{f_2(x)} f_{2,6}(x) - Af_{3,6}(x) - Bf_{4,6}(x) - Df_{5,6}(x) \right) \in C_6 \\ \Rightarrow & f_6(x) \mid \frac{x^n-1}{f_1(x)} \left(f_{1,6}(x) - \frac{f_{1,2}(x)}{f_2(x)} f_{2,6}(x) - Af_{3,6}(x) - Bf_{4,6}(x) - Df_{5,6}(x) \right), \end{aligned}$$

where,

$$\begin{aligned} A &= \frac{x^n-1}{f_1(x)} \frac{f_{1,2}(x)}{f_2(x)}, \\ B &= \frac{x^n-1}{f_1(x)} \left(\frac{f_{1,3}(x) - \frac{f_{1,2}(x)}{f_2(x)} f_{2,3}(x)}{f_3(x)} \right), \\ C &= \frac{x^n-1}{f_1(x)} \left(\frac{-f_{1,4}(x) + \frac{f_{1,2}(x)}{f_2(x)} f_{2,4}(x) + Af_{3,4}(x)}{f_4(x)} \right) \text{ and} \\ D &= \frac{x^n-1}{f_1(x)} \left(\frac{-f_{1,5}(x) + \frac{f_{1,2}(x)}{f_2(x)} f_{2,5}(x) + Af_{3,5}(x) + Bf_{4,5}(x)}{f_5(x)} \right). \end{aligned}$$

Similarly, we get the results for rest of the values of i .

$$(9) \text{ For } i = 7, \text{ we have } \frac{x^n-1}{f_1(x)} \left(A_1 - AA_2 - BA_3 - DA_4 - EA_5 - FA_6 \right) \in C.$$

Since

$$vw \frac{x^n-1}{f_1(x)} \left(f_{1,7}(x) - Af_{2,7}(x) - Bf_{3,7}(x) - Df_{4,7}(x) - Ef_{5,7}(x) - Ff_{6,7}(x) \right) \in C \pmod{uvw}, \text{ therefore,}$$

$$\begin{aligned} & \frac{x^n-1}{f_1(x)} \left(f_{1,7}(x) - Af_{2,7}(x) - Bf_{3,7}(x) - Df_{4,7}(x) - Ef_{5,7}(x) - Ff_{6,7}(x) \right) \in C_7 \\ \Rightarrow & f_7(x) \mid \frac{x^n-1}{f_1(x)} \left(f_{1,7}(x) - Af_{2,7}(x) - Bf_{3,7}(x) - Df_{4,7}(x) - Ef_{5,7}(x) - Ff_{6,7}(x) \right), \end{aligned}$$

where

$$\begin{aligned} A &= \left(\frac{f_{1,2}(x)}{f_2(x)} \right), \\ B &= \left(f_{1,3}(x) - A \frac{f_{2,3}(x)}{f_3(x)} \right), \\ D &= \left(\frac{f_{1,4}(x) - Af_{2,4}(x) - Bf_{3,4}(x)}{f_4(x)} \right), \end{aligned}$$

$$E = \left(\frac{f_{1,5}(x) - Af_{2,5}(x) - Bf_{3,5}(x) - Df_{4,5}(x)}{f_5(x)} \right) \text{ and}$$

$$F = \left(\frac{f_{1,6}(x) - Af_{2,6}(x) - Bf_{3,6}(x) - Df_{4,6}(x) - Ef_{5,6}(x)}{f_6(x)} \right).$$

Similarly, we get the results for rest of the values of i .

(10) We have

$$\begin{aligned} & \frac{x^n-1}{f_1(x)} \left(A_1 - AA_2 - BA_3 - DA_4 - EA_5 - FA_6 - GA_7 \right) \\ &= uvw \frac{x^n-1}{f_1(x)} * \left(f_{1,8}(x) - Af_{2,8}(x) - Bf_{3,8}(x) - Df_{4,8}(x) - Ef_{5,8}(x) \right. \\ & \quad \left. - Ff_{6,8}(x) - Gf_{7,8}(x) \right) \in C. \end{aligned}$$

Therefore,

$$\begin{aligned} & \frac{x^n-1}{f_1(x)} \left(f_{1,8}(x) - Af_{2,8}(x) - Bf_{3,8}(x) - Df_{4,8}(x) - Ef_{5,8}(x) \right. \\ & \quad \left. - Ff_{6,8}(x) - Gf_{7,8}(x) \right) \in C_8 \\ & \Rightarrow f_8(x) \mid \frac{x^n-1}{f_1(x)} * \left(f_{1,8}(x) - Af_{2,8}(x) - Bf_{3,8}(x) - Df_{4,8}(x) - Ef_{5,8}(x) \right. \\ & \quad \left. - Ff_{6,8}(x) - Gf_{7,8}(x) \right), \end{aligned}$$

where

$$\begin{aligned} A &= \left(\frac{f_{1,2}(x)}{f_2(x)} \right), \\ B &= \left(f_{1,3}(x) - A \frac{f_{2,3}(x)}{f_3(x)} \right), \\ D &= \left(\frac{f_{1,4}(x) - Af_{2,4}(x) - Bf_{3,4}(x)}{f_4(x)} \right), \\ E &= \left(\frac{f_{1,5}(x) - Af_{2,5}(x) - Bf_{3,5}(x) - Df_{4,5}(x)}{f_5(x)} \right), \\ F &= \left(\frac{f_{1,6}(x) - Af_{2,6}(x) - Bf_{3,6}(x) - Df_{4,6}(x) - Ef_{5,6}(x)}{f_6(x)} \right) \text{ and} \\ G &= \left(\frac{f_{1,7}(x) - Af_{2,7}(x) - Bf_{3,7}(x) - Df_{4,7}(x) - Ef_{5,7}(x) - Ff_{6,7}(x)}{f_7(x)} \right). \end{aligned}$$

(11) For $i = 4$, we have $uA_3 = uvf_{2,3}(x) + wvf_{2,5}(x) + uvwf_{2,7}(x) \in C$. Therefore, $uvf_{2,3}(x) \in C \pmod w \Rightarrow f_{2,3}(x) \in C_4 \Rightarrow f_4(x) \mid f_{2,3}(x)$. Similarly we get the results for rest of the values of i .

(12) For $i = 4$, we have $(vA_1 - \frac{f_1(x)}{f_3(x)}A_3) \in C$. Since $uv \left(f_{1,2}(x) - \frac{f_1(x)}{f_3(x)}f_{3,4}(x) \right) \in C \pmod w$, therefore,

$$\left(f_{1,2}(x) - \frac{f_1(x)}{f_3(x)}f_{3,4}(x) \right) \in C_4 \Rightarrow f_4(x) \mid \left(f_{1,2}(x) - \frac{f_1(x)}{f_3(x)}f_{3,4}(x) \right).$$

Similarly, we get the results for rest of the values of i .

(13) For $i = 7$, we have $\left(wA_2 - \frac{f_2(x)}{f_6(x)}A_6\right) \in C$. Since $vw\left(f_{2,3}(x) - \frac{f_2(x)}{f_6(x)}f_{6,7}(x)\right) \in C \pmod w$, therefore, $f_7(x) \mid \left(f_{2,3}(x) - \frac{f_2(x)}{f_6(x)}f_{6,7}(x)\right)$. Similarly we get the result for $i = 8$.

(14) For $i = 7$, we have $\left(wA_1 - \frac{f_1(x)}{f_5(x)}A_5 - \frac{f_{1,2}(x) - \frac{f_1(x)}{f_5(x)}f_{5,6}(x)}{f_6(x)}A_6\right) \in C$. Since $vw\left(f_{1,3}(x) - \frac{f_1(x)}{f_5(x)}f_{5,7}(x) - \frac{f_{1,2}(x) - \frac{f_1(x)}{f_5(x)}f_{5,6}(x)}{f_6(x)}f_{6,7}(x)\right) \in C \pmod{uvw}$, therefore,

$$\begin{aligned} & \left(f_{1,3}(x) - \frac{f_1(x)}{f_5(x)}f_{5,7}(x) - \frac{f_{1,2}(x) - \frac{f_1(x)}{f_5(x)}f_{5,6}(x)}{f_6(x)}f_{6,7}(x)\right) \in C_7, \\ \Rightarrow & f_7(x) \mid \left(f_{1,3}(x) - \frac{f_1(x)}{f_5(x)}f_{5,7}(x) - \frac{f_{1,2}(x) - \frac{f_1(x)}{f_5(x)}f_{5,6}(x)}{f_6(x)}f_{6,7}(x)\right). \end{aligned}$$

Similarly, we get the results for rest of the values of i .

(15) We have $vA_4 \in C$. Since $vwf_{4,5}(x) \in C \pmod{uvw}$, therefore, $f_{4,5}(x) \in C_7 \Rightarrow f_7(x) \mid f_{4,5}(x)$. Similarly by taking vA_3 we can show $f_7(x) \mid f_{3,5}(x)$.

(16) We have $uvA_2 = uvwf_{2,5}(x) \in C$. Therefore,

$$f_{2,5}(x) \in C_8 \Rightarrow f_8(x) \mid f_{2,5}(x).$$

(17) We have $vA_3 - \frac{f_{3,5}(x)}{f_7(x)}A_7 = uvw\left(f_{3,6}(x) - \frac{f_{3,5}(x)}{f_7(x)}f_{7,8}(x)\right) \in C_8$. Therefore,

$$\left(f_{3,6}(x) - \frac{f_{3,5}(x)}{f_7(x)}f_{7,8}(x)\right) \in C_8 \Rightarrow f_8(x) \mid \left(f_{3,6}(x) - \frac{f_{3,5}(x)}{f_7(x)}f_{7,8}(x)\right).$$

(18) We have $vA_4 - \frac{f_{4,5}(x)}{f_7(x)}A_7 = uvw\left(f_{4,6}(x) - \frac{f_{4,5}(x)}{f_7(x)}f_{7,8}(x)\right) \in C$. Therefore,

$$\left(f_{4,6}(x) - \frac{f_{4,5}(x)}{f_7(x)}f_{7,8}(x)\right) \in C_8 \Rightarrow f_8(x) \mid \left(f_{4,6}(x) - \frac{f_{4,5}(x)}{f_7(x)}f_{7,8}(x)\right).$$

(19) We have $vA_5 - \frac{f_5(x)}{f_7(x)}A_7 = uvw\left(f_{5,6}(x) - \frac{f_5(x)}{f_7(x)}f_{7,8}(x)\right) \in C$. Therefore,

$$\left(f_{5,6}(x) - \frac{f_5(x)}{f_7(x)}f_{7,8}(x)\right) \in C_8 \Rightarrow f_8(x) \mid \left(f_{5,6}(x) - \frac{f_5(x)}{f_7(x)}f_{7,8}(x)\right).$$

(20) We have $wA_1 - \frac{f_1(x)}{f_5(x)}A_5 - AA_6 - BA_7 = uvw\left(f_{1,4}(x) - \frac{f_1(x)}{f_5(x)}f_{5,8}(x) - Af_{6,8}(x) - Bf_{7,8}(x)\right) \in C$. Therefore,

$$\begin{aligned} & \left(f_{1,4}(x) - \frac{f_1(x)}{f_5(x)}f_{5,8}(x) - Af_{6,8}(x) - Bf_{7,8}(x)\right) \in C_8 \\ \Rightarrow & f_8(x) \mid \left(f_{1,4}(x) - \frac{f_1(x)}{f_5(x)}f_{5,8}(x) - Af_{6,8}(x) - Bf_{7,8}(x)\right), \end{aligned}$$

where $A = \left(\frac{f_{1,2}(x) - \frac{f_1(x)}{f_5(x)}f_{5,6}(x)}{f_6(x)}\right)$ and $B = \left(\frac{f_{1,3}(x) - \frac{f_1(x)}{f_5(x)}f_{5,7}(x) - Af_{6,7}(x)}{f_6(x)}\right)$. \square

Theorem 3. *If $C = \langle A_1, A_2, \dots, A_8 \rangle$ is a cyclic code over the ring $R_{u^2, v^2, w^2, p}$ then C is a free cyclic code if and only if $f_1(x) = f_8(x)$. In this case, we have $C = \langle A_1 \rangle$ and $A_1 | (x^n - 1)$ in $R_{u^2, v^2, w^2, p}[x]$.*

Proof. Let $f_1(x) = f_8(x)$. Since $f_8(x) | f_4(x) | f_2(x) | f_1(x)$, $f_4(x) | f_3(x) | f_1(x)$, $f_6(x) | f_2(x) | f_1(x)$, $f_8(x) | f_6(x) | f_5(x) | f_1(x)$, $f_8(x) | f_7(x) | f_3(x) | f_1(x)$ and $f_7(x) | f_5(x) | f_1(x)$, therefore, we get $f_1(x) = f_2(x) = f_3(x) = f_4(x) = f_5(x) = f_6(x) = f_7(x) = f_8(x)$. Let $B_1 = f_1(x) + uf_{1,2}(x) + vf_{1,3}(x) + uvf_{1,4}(x)$, $B_2 = uf_2(x) + vf_{2,3}(x) + uvf_{2,4}(x)$, $B_3 = vf_3(x) + uvf_{3,4}(x)$, $B_4 = uvf_4(x)$. Then we have $\text{Im}\phi = \langle B_1, B_2, B_3, B_4 \rangle$ and $\text{Ker}\phi = w\langle B_5, B_6, B_7, B_8 \rangle$, where $B_5 = f_5(x) + uf_{5,6}(x) + vf_{5,7}(x) + uvf_{5,8}(x)$, $B_6 = uf_6(x) + vf_{6,7}(x) + uvf_{6,8}(x)$, $B_7 = vf_7(x) + uvf_{7,8}(x)$, $B_8 = uvf_8(x)$. From [8, Proposition 3.3], we get $\text{Im}\phi = \langle B_1 \rangle$ and $\text{Ker}\phi = w\langle B_5 \rangle$. Therefore, we have $C = \langle f_1(x) + uf_{1,2}(x) + vf_{1,3}(x) + uvf_{1,4}(x) + wf_{1,5}(x) + uvf_{1,6}(x) + vwf_{1,7}(x) + uvwf_{1,8}(x), wf_5(x) + uvf_{5,6}(x) + vwf_{5,7}(x) + uvwf_{5,8}(x) \rangle$. Now to show $C = \langle A_1 \rangle$ we show that $f_{1,2}(x) = f_{5,6}(x)$, $f_{1,3}(x) = f_{5,7}(x)$ and $f_{1,4}(x) = f_{5,8}(x)$. Since $wA_1 - A_5 = uv(f_{1,2}(x) - f_{5,6}(x)) + vw(f_{1,3}(x) - f_{5,7}(x)) + uvw(f_{1,4}(x) - f_{5,8}(x)) \in C$, this gives, $(f_{1,2}(x) - f_{5,6}(x)) \in C_6 = \langle f_6(x) \rangle$. Therefore, $f_6(x) | (f_{1,2}(x) - f_{5,6}(x))$. Since $(\deg(f_{1,2}(x)), \deg(f_{5,6}(x))) < \deg(f_6(x))$, this implies that $f_{1,2}(x) - f_{5,6}(x) = 0$. Thus, $f_{1,2}(x) = f_{5,6}(x)$. Hence, we get $wA_1 - A_5 = vw(f_{1,3}(x) - f_{5,7}(x)) + uvw(f_{1,4}(x) - f_{5,8}(x)) \in C$. This gives, $(f_{1,3}(x) - f_{5,7}(x)) \in C_7 = \langle f_7(x) \rangle$. Therefore, $f_7(x) | (f_{1,3}(x) - f_{5,7}(x))$. Since $(\deg(f_{1,3}(x)), \deg(f_{5,7}(x))) < \deg(f_7(x))$, this implies that $(f_{1,3}(x) - f_{5,7}(x)) = 0$. Therefore, $f_{1,3}(x) = f_{5,7}(x)$. Finally, we have $wA_1 - A_5 = uvw(f_{1,4}(x) - f_{5,8}(x)) \in C$. This gives, $(f_{1,4}(x) - f_{5,8}(x)) \in C_8 = \langle f_8(x) \rangle$. Therefore, $f_8(x) | (f_{1,4}(x) - f_{5,8}(x))$. Since

$$(\deg(f_{1,4}(x)), \deg(f_{5,8}(x))) < \deg(f_8(x)),$$

this implies that $(f_{1,4}(x) - f_{5,8}(x)) = 0$. Therefore, $f_{1,4}(x) = f_{5,8}(x)$. This shows that $wA_1 = A_5$. Hence $C = \langle A_1 \rangle$ and $C \simeq R_{u^2, v^2, w^2, p}^{n - \deg(f_1(x))}$. Conversely, if C is a free cyclic code, we must have $C = \langle A_1 \rangle$. Since $uvwf_8(x) \in C$, we have $uvwf_8(x) = uvw\alpha f_1(x)$ for some $\alpha \in \mathbb{F}_p$. Note that $f_8(x) | f_1(x)$, hence by comparing the coefficients both sides, we get $f_1(x) = f_8(x)$. Now for the second condition, by the division algorithm, we have $x^n - 1 = A_1q(x) + r(x)$, where $r(x) = 0$ or $\deg(r(x)) < \deg(f_1(x))$. This implies that $r(x) = (x^n - 1) - A_1q(x) \in C$. Since A_1 is the lowest degree polynomial in C , so $r(x) = 0$. Hence, $A_1 | (x^n - 1)$ in $R_{u^2, v^2, w^2, p}[x]$. \square

2.1. When $\gcd(n, p) = 1$

Let n be a positive integer relatively prime to p . First, we slightly refine Theorem 3.4 of [8], which gives the structure of a cyclic code over the ring $R_{u^2, v^2, p} = \mathbb{F}_p[u, v] / \langle u^2, v^2, uv - vu \rangle$. Let $C_{u, v}$ be a cyclic code over the ring $R_{u^2, v^2, p}$. From Theorem 3.4 of [8], we have $C_{u, v} = \langle g(x) + ua_1(x) + uvr_1(x), va_2(x) + uvva_3(x) \rangle$ with $a_3 | a_1(x) | g(x) | (x^n - 1)$ and $a_3(x) | a_2(x) | g(x) | (x^n - 1)$. Thus, $\frac{x^n - 1}{g(x)} \frac{x^n - 1}{a_1(x)} (g(x) + ua_1(x) + uvr_1(x)) = uv \frac{x^n - 1}{g(x)} \frac{x^n - 1}{a_1(x)} r_1(x) \in C_{u, v}$. This gives, $\frac{x^n - 1}{g(x)} \frac{x^n - 1}{a_1(x)} r_1(x) \in$

$C_4 = \text{Tor}(\text{Tor}(C_{u,v})) = \langle a_3(x) \rangle$ (see Page 165 of [8]). Hence, $a_3(x) \mid \frac{x^n-1}{g(x)} \frac{x^n-1}{a_1(x)} r_1(x)$. Since n is relatively prime to p , $x^n - 1$ can be uniquely factored as product of distinct irreducible factors. Therefore, we must have $\text{g.c.d.}(a_3(x), \frac{x^n-1}{g(x)}) = \text{g.c.d.}(a_3(x), \frac{x^n-1}{a_1(x)}) = 1$. This gives, $a_3(x) \mid r_1(x)$. But, from Theorem 3.1 of [8], we have $\deg(r_1(x)) < \deg(a_3(x))$. This gives, $r_1(x) = 0$. Thus we have proved the following theorem.

Theorem 4. *Let $C_{u,v}$ be a cyclic code over the ring $R_{u^2, v^2, p}$ of length n . If n is relatively prime to p , then we have $C_{u,v} = \langle f_1(x) + uf_2(x), vf_3(x) + uvf_4(x) \rangle$ with $f_4(x) \mid f_2(x) \mid f_1(x) \mid (x^n - 1)$ and $f_4(x) \mid f_3(x) \mid f_1(x) \mid (x^n - 1)$.*

If $C = \langle A_1, A_2, \dots, A_8 \rangle$ is a cyclic code of length n over the ring $R_{u^2, v^2, w^2, p}$ then we have $\text{Im}\phi = \langle A_1, A_2, A_3, A_4 \rangle$ and $\ker\phi = w\langle A_5, A_6, A_7, A_8 \rangle$. (See Equation 2.1 for the definition of ϕ). Note that we can consider $\text{Im}\phi$ and $\ker\phi$ as cyclic codes over the ring $R_{u^2, v^2, p}$. Since n is relatively prime to p , from the above theorem, we have $\text{Im}\phi = \langle f_1(x) + uf_2(x), vf_3(x) + uvf_4(x) \rangle$ and $\ker\phi = w\langle f_5(x) + uf_6(x), vf_7(x) + uvf_8(x) \rangle$ with $f_4(x) \mid f_2(x) \mid f_1(x) \mid (x^n - 1)$, $f_4(x) \mid f_3(x) \mid f_1(x) \mid (x^n - 1)$, $f_8(x) \mid f_6(x) \mid f_5(x) \mid (x^n - 1)$ and $f_8(x) \mid f_7(x) \mid f_5(x) \mid (x^n - 1)$. We also have the conditions $f_5(x) \mid f_1(x)$, $f_6(x) \mid f_2(x)$ and $f_7(x) \mid f_3(x)$ (from Conditions 1 and 2 of Theorem 2). Therefore, the code C can be written as $C = \langle f_1(x) + uf_2(x) + w(f_{1,5}(x) + uf_{1,6}(x) + vf_{1,7}(x) + uvf_{1,8}(x)), f_3(x) + uf_4(x) + w(f_{3,5}(x) + uf_{3,6}(x) + vf_{3,7}(x) + uvf_{3,8}(x)), w(f_5(x) + uf_6(x)), w(vf_7(x) + uvf_8(x)) \rangle$ with the same conditions as above on $f_i(x)$'s. From Condition 4 of Theorem 2, for $i = 1$ and $4 \leq j \leq 7$, we get $f_{1+j}(x) \mid \frac{x^n-1}{f_1(x)} \frac{x^n-1}{f_2(x)} \cdots \frac{x^n-1}{f_j(x)} f_{1,1+j}(x)$. Since n is relatively prime to p , $x^n - 1$ can be uniquely factored as product of distinct irreducible factors. Therefore, we must have $\text{g.c.d.}(f_{1+j}(x), \frac{x^n-1}{f_k(x)}) = 1$, for $1 \leq k \leq j$. This gives $f_{1+j}(x) \mid f_{1,1+j}(x)$. From Theorem 1, we have $\deg(f_{1,j+1}(x)) < \deg(f_{1+j}(x))$, for $4 \leq j \leq 7$. This gives $f_{1,1+j}(x) = 0$ for $4 \leq j \leq 7$. Similarly, from Condition 4 of Theorem 2, for $i = 3$ and $2 \leq j \leq 5$, we can show that $f_{3,3+j}(x) = 0$. Thus we have proved the following theorem.

Theorem 5. *Let $C = \langle A_1, A_2, \dots, A_8 \rangle$ be a cyclic code over the ring $R_{u^2, v^2, w^2, p}$ of length n . If n is relatively prime to p , then we have*

$C = \langle f_1(x) + uf_2(x), vf_3(x) + uvf_4(x), w(f_5(x) + uf_6(x)), w(vf_7(x) + uvf_8(x)) \rangle$
with the conditions:

$$f_4(x) \mid f_2(x) \mid f_1(x) \mid (x^n - 1), f_4(x) \mid f_3(x) \mid f_1(x), f_8(x) \mid f_6(x) \mid f_5(x) \mid (x^n - 1), \\ f_8(x) \mid f_7(x) \mid f_5(x) \mid f_1(x), f_6(x) \mid f_2(x) \text{ and } f_7(x) \mid f_3(x).$$

3. Ranks and minimal spanning sets

We follow Dougherty and Shiromoto [7, page 401] for the definition of the rank of a code C . We first prove the number of lemmas that we use to find the rank and the minimal spanning set of cyclic codes over $R_{u^2, v^2, w^2, p}$.

Lemma 1. *Let C be a cyclic code over the ring $R_{u^2, v^2, w^2, p}$. If $C = \langle A_1, A_2, \dots, A_8 \rangle$, then polynomials in C in the following forms can be written as follows:*

- (1) $w(p_0(x) + up_1(x) + vp_2(x) + uvp_3(x)) = q_5(x)A_5 + q_6(x)A_6 + q_7(x)A_7 + q_8(x)A_8$,
- (2) $w(up_1(x) + vp_2(x) + uvp_3(x)) = q_6(x)A_6 + q_7(x)A_7 + q_8(x)A_8$,
- (3) $w(vp_2(x) + uvp_3(x)) = q_7(x)A_7 + q_8(x)A_8$,
- (4) $w(uvp_3(x)) = q_8(x)A_8$,

for some $q_i(x) \in \mathbb{F}_p[x]$, $5 \leq i \leq 8$.

Proof. (1) Let $A' = w(p_0(x) + up_1(x) + vp_2(x) + uvp_3(x)) \in C$. Thus, $p_0(x) \in C_5 = \langle f_5(x) \rangle$. This gives, $p_0(x) = q_5(x)f_5(x)$ for some $q_5(x) \in \mathbb{F}_p[x]$. Therefore, $A' - q_5(x)A_5 = w((p_1(x) - q_5(x)f_{5,6}(x)) + u(p_2(x) - q_5(x)f_{5,7}(x)) + uv(p_3(x) - q_5(x)f_{5,8}(x))) \in C$. Thus, $p_1(x) - q_5(x)f_{5,6}(x) \in C_6 = \langle f_6(x) \rangle$. Therefore, $(p_1(x) - q_5(x)f_{5,6}(x)) = q_6(x)f_6(x)$ for some $q_6(x) \in \mathbb{F}_p[x]$. Again,

$$\begin{aligned} A' - q_5(x)A_5 - q_6(x)A_6 &= w(v(p_2(x) - q_5(x)f_{5,7}(x) - q_1(x)f_{6,7}(x)) \\ &\quad + uv(p_3(x) - q_5(x)f_{5,8}(x) - q_1(x)f_{6,8}(x))) \in C. \end{aligned}$$

Thus, $(p_2(x) - q_5(x)f_{5,7}(x) - q_1(x)f_{6,7}(x)) \in C_7 = \langle f_7(x) \rangle$. Therefore, $(p_2(x) - q_5(x)f_{5,7}(x) - q_6(x)f_{6,7}(x)) = q_7(x)f_7(x)$ for some $q_7(x) \in \mathbb{F}_p[x]$. Again,

$$\begin{aligned} A' - q_5(x)A_5 - q_6(x)A_6 - q_7(x)A_7 \\ = w(uv(p_3(x) - q_5(x)f_{5,8}(x) - q_6(x)f_{6,8}(x) - q_7(x)f_{7,8}(x))) \in C. \end{aligned}$$

Thus, $(p_3(x) - q_5(x)f_{5,8}(x) - q_6(x)f_{6,8}(x) - q_7(x)f_{7,8}(x)) \in C_8 = \langle f_8(x) \rangle$. Therefore, $(p_3(x) - q_5(x)f_{5,8}(x) - q_6(x)f_{6,8}(x) - q_7(x)f_{7,8}(x)) = q_8(x)f_8(x)$ for some $q_8(x) \in \mathbb{F}_p[x]$. That is, $A' - q_5(x)A_5 - q_6(x)A_6 - q_7(x)A_7 - q_8(x)A_8 = 0 \Rightarrow A' = q_5(x)A_5 + q_6(x)A_6 + q_7(x)A_7 + q_8(x)A_8$. This proves Statement (1). The proof of other cases are similar to the proof of Statement (1). \square

Lemma 2. *Let C be a cyclic code over the ring $R_{u^2, v^2, w^2, p}$. If $C = \langle A_1, A_2, \dots, A_8 \rangle$ and $\deg(f_i(x)) = t_i$, $1 \leq i \leq 8$, then the following conditions hold:*

- (1) $x^{t_i - t_8} A_8 = c_i u_{8-i} A_i + q_8(x) A_8$, $1 \leq i \leq 7$, where $\deg(q_8(x)) < t_i - t_8$, $u_1 = u, u_2 = v, u_3 = uv, u_4 = w, u_5 = uw, u_6 = vw$ and $u_7 = uvw$,
- (2) $x^{t_5 - t_7} A_7 = c_5 v A_5 - q'_7(x) A_7 - q'_8(x) A_8$, where $\deg(q'_7(x)) < t_5 - t_7$,
- (3) $x^{t_3 - t_7} A_7 = c_3 w A_3 - q'_7(x) A_7 - q'_8(x) A_8$, where $\deg(q'_7(x)) < t_3 - t_7$,
- (4) $x^{t_1 - t_7} A_7 = c_1 v w A_1 - q'_7(x) A_7 - q'_8(x) A_8$, where $\deg(q'_7(x)) < t_1 - t_7$,
- (5) $x^{t_5 - t_6} A_6 = c_5 u A_5 - q'_6(x) A_6 - q'_7(x) A_7 - q'_8(x) A_8$, where $\deg(q'_6(x)) < t_5 - t_6$,
- (6) $x^{t_2 - t_6} A_6 = c_2 w A_2 - q'_6(x) A_6 - q'_7(x) A_7 - q'_8(x) A_8$, where $\deg(q'_6(x)) < t_2 - t_6$,
- (7) $x^{t_1 - t_6} A_6 = c_1 u w A_1 - q'_6(x) A_6 - q'_7(x) A_7 - q'_8(x) A_8$, where $\deg(q'_6(x)) < t_1 - t_6$, and
- (8) $x^{t_1 - t_5} A_5 = c_1 w A_1 + q_5(x) A_5 + q_6(x) A_6 + q_7(x) A_7 + q_8(x) A_8$, where $\deg(q_5(x)) < t_1 - t_5$,

$c_i \in \mathbb{F}_p$ and $q_i(x), q'_i(x) \in \mathbb{F}_p[x]$.

Proof. (1) From Condition (1) of Theorem 2, we have $f_8(x)|f_i(x)$, $1 \leq i \leq 7$. Thus, $f_i(x) = s_i(x)f_8(x)$ for some $s_i(x) \in \mathbb{F}_p[x]$. This can be written as, $f_i(x) = (s_{i0} + xs_{i1} + \cdots + x^{t_i-t_8}s_{i(t_i-t_8)})f_8(x)$, where $s_{ij} \in \mathbb{F}_p$. Clearly $s_{i(t_i-t_8)} \neq 0$. Therefore, $u_{8-i}A_i - s_i(x)A_8 = uvw(f_i(x) - s_i(x)f_8(x)) = 0$. This gives, $x^{t_i-t_8}A_8 = s_{i(t_i-t_8)}^{-1}u_{8-i}A_i - s_{i(t_i-t_8)}^{-1}(s_{i0} + xs_{i1} + \cdots + x^{t_i-t_8-1}s_{i(t_i-t_8-1)})A_8$. Hence, $x^{t_i-t_8}A_8 = c_i u_{8-i}A_i + q_8(x)A_8$, where $\deg(q_8(x)) < t_i - t_8$.

(2) From Condition (2) of Theorem 2, we have $f_7(x)|f_5(x)$. Thus, $f_5(x) = s_5(x)f_7(x)$ for some $s_5(x) \in F_p[x]$. This can be written as $f_5(x) = (s_{50} + xs_{51} + \cdots + x^{t_5-t_7}s_{5(t_5-t_7)})f_7(x)$. This together with Condition (4) of Lemma 1, we get $vA_5 - s_5(x)A_7 = w(uv(f_{5,6}(x) - s_5(x)f_{7,8}(x))) = q_8(x)A_8$. Thus,

$$(3.1) \quad s_5(x)A_7(x) = vA_5 - q_8(x)A_8.$$

This can be written as $x^{t_5-t_7}A_7 = s_{5(t_5-t_7)}^{-1}vA_5 - s_{5(t_5-t_7)}^{-1}(s_{50} + xs_{51} + \cdots + x^{t_5-t_7-1}s_{5(t_5-t_7-1)})A_7(x) - s_{5(t_5-t_7)}^{-1}q_8(x)A_8$. Thus,

$$(3.2) \quad x^{t_5-t_7}A_7 = c_5vA_5 - q_7'(x)A_7 - q_8'(x)A_8,$$

where $\deg(q_7'(x)) = t_5 - t_7 - 1 < t_5 - t_7$.

(3) The proof is similar to Condition 2.

(4) The proof is similar to Condition 2.

(5) From Condition (2) of Theorem 2, we have $f_6(x)|f_5(x)$. Thus, $f_5(x) = s_5(x)f_6(x)$ for some $s_5(x) \in F_p[x]$. This can be written as, $f_5(x) = (s_{50} + s_{51}x + \cdots + s_{5(t_5-t_6)}x^{t_5-t_6})f_6(x)$, where $s_{5i} \in \mathbb{F}_p$. This together with Condition (3) of Lemma 1, we get $uA_5 - s_5(x)A_6(x) = w(v(-s_5(x)f_{6,7}(x)) + uv(f_{5,7}(x) - s_5(x)f_{6,8}(x))) = q_7(x)A_7 + q_8(x)A_8 \in C$. Thus,

$$(3.3) \quad s_5(x)A_6(x) = uA_5 - q_7(x)A_7 - q_8(x)A_8.$$

This can be written as $x^{t_5-t_6}A_6 = s_{5(t_5-t_6)}^{-1}uA_5 - s_{5(t_5-t_6)}^{-1}(s_{50} + xs_{51} + \cdots + x^{t_5-t_6-1}s_{5(t_5-t_6-1)})A_6 - s_{5(t_5-t_6)}^{-1}q_7(x)A_7 - s_{5(t_5-t_6)}^{-1}q_8(x)A_8$, i.e.,

$$(3.4) \quad x^{t_5-t_6}A_6 = c_5uA_5 - q_6'(x)A_6 - q_7'(x)A_7 - q_8'(x)A_8,$$

where $\deg(q_6'(x)) = t_5 - t_6 - 1 < t_5 - t_6$.

(6) The proof is similar to Condition 5.

(7) The proof is similar to Condition 5.

(8) By the division algorithm, we have

$$(3.5) \quad \begin{aligned} & x^{t_1-t_5}(f_5(x) + uf_{5,6}(x) + vf_{5,7}(x) + wvf_{5,8}(x)) \\ &= c_1A_1 + (p_0(x) + up_1(x) + vp_2(x) + wvp_3(x) + wp_4(x) + wwp_5(x) \\ & \quad + vwp_6(x) + uvwp_7(x)), \end{aligned}$$

where $\deg(p_0(x)) < \deg(f_1(x)) = t_1$. Multiplying Eq. (3.5) by w and applying Condition 1 of Lemma 1 gives, $x^{t_1-t_5}A_5 - c_1wA_1 = w(p_0(x) + up_1(x) + vp_2(x) + wvp_3(x)) = q_5(x)A_5 + q_6(x)A_6 + q_7(x)A_7 + q_8(x)A_8$. That is,

$$(3.6) \quad x^{t_1-t_5}A_5 = c_1wA_1 + q_5(x)A_5 + q_6(x)A_6 + q_7(x)A_7 + q_8(x)A_8.$$

We have $p_0(x) = q_5(x)f_5(x)$, thus, $\deg(q_5(x)) + \deg(f_5(x)) = \deg(p_0(x)) < t_1$. Hence, $\deg(q_5(x)) < t_1 - t_5$. \square

Theorem 6. *Let C be a cyclic code of length n over $R_{u^2, v^2, w^2, p}$. If $C = \langle A_1, A_2, A_3, A_4, A_5, A_6, A_7, A_8 \rangle$ with $t_i = \deg(f_i(x))$, $1 \leq i \leq 8$, $t'_4 = \min\{t_2, t_3\}$, $t'_6 = \min\{t_2, t_5\}$, $t'_7 = \min\{t_3, t_5\}$ and $t'_8 = \min\{t_4, t_6, t_7\}$, then C has rank $n + 2t_1 + t'_4 + t'_6 + t'_7 + t'_8 - t_2 - t_3 - t_4 - t_5 - t_6 - t_7 - t_8$. The minimal spanning set B of the code C is $B = \{A_1, xA_1, \dots, x^{n-t_1-1}A_1, A_2, xA_2, \dots, x^{t_1-t_2-1}A_2, A_3, xA_3, \dots, x^{t_1-t_3-1}A_3, A_4, xA_4, \dots, x^{t'_4-t_4-1}A_4, A_5, xA_5, \dots, x^{t_1-t_5-1}A_5, A_6, xA_6, \dots, x^{t'_6-t_6-1}A_6, A_7, xA_7, \dots, x^{t'_7-t_7-1}A_7, A_8, xA_8, \dots, x^{t'_8-t_8-1}A_8\}$.*

Proof. It suffices to show that B spans the set $B' = \{A_1, xA_1, \dots, x^{n-t_1-1}A_1, A_2, xA_2, \dots, x^{n-t_2-1}A_2, A_3, xA_3, \dots, x^{n-t_3-1}A_3, A_4, xA_4, \dots, x^{n-t_4-1}A_4, A_5, xA_5, \dots, x^{n-t_5-1}A_5, A_6, xA_6, \dots, x^{n-t_6-1}A_6, A_7, xA_7, \dots, x^{n-t_7-1}A_7, A_8, xA_8, \dots, x^{n-t_8-1}A_8\}$. To show B spans B' , we write the set B' as $B' = B_1 \cup B_2$, where $B_1 = \{A_1, xA_1, \dots, x^{n-t_1-1}A_1, A_2, xA_2, \dots, x^{n-t_2-1}A_2, A_3, xA_3, \dots, x^{n-t_3-1}A_3, A_4, xA_4, \dots, x^{n-t_4-1}A_4\}$ and $B_2 = \{A_5, xA_5, \dots, x^{n-t_5-1}A_5, A_6, xA_6, \dots, x^{n-t_6-1}A_6, A_7, xA_7, \dots, x^{n-t_7-1}A_7, A_8, xA_8, \dots, x^{n-t_8-1}A_8\}$. First we show that B spans B_2 and then we show that B spans B_1 . To show B spans B_2 we divide the proof in twelve cases.

Case (1). Let $t'_8 = t_7$, $t'_7 = t_5$ and $t'_6 = t_5$. We first show that the element $x^{t_7-t_8}A_8 \in B_2 - B$ is linear combinations of some elements of B and then we show that other elements of the set $B_2 - B$ are linear combinations of elements of B . From Statement 1 of Lemma 2,

$$(3.7) \quad x^{t_7-t_8}A_8 = c_7uA_7 + q_8(x)A_8,$$

where $\deg(q_8(x)) < t_7 - t_8$. Therefore, $x^{t_7-t_8}A_8 \in \text{Span}(B)$. Multiplying Equation (3.7) by $x, x^2, x^3, \dots, x^{t_5-t_7-1}$ and then putting the value of $x^{t_7-t_8}A_8$ in the equation obtained, we can show that $x^{t_7-t_8+1}A_8, x^{t_7-t_8+2}A_8, \dots, x^{t_5-t_8-1}A_8 \in \text{Span}(B)$. From Statement 1 of Lemma 2, we have

$$(3.8) \quad x^{t_5-t_8}A_8 = c_5uvA_5 + q_8(x)A_8,$$

where $\deg(q_8(x)) < t_5 - t_8$. Therefore, $x^{t_5-t_8}A_8 \in \text{Span}(B)$. Arguing as above, we can show that the terms $x^{t_5-t_8+1}A_8, x^{t_5-t_8+2}A_8, \dots, x^{t_1-t_8-1}A_8 \in \text{Span}(B)$. Again, from Statement 1 of Lemma 2, we have

$$(3.9) \quad x^{t_1-t_8}A_8 = c_5uvwA_1 + q_8(x)A_8,$$

where $\deg(q_8(x)) < t_1 - t_8$. As above, we can show that $x^{t_1-t_8}A_8, x^{t_1-t_8+1}A_8, \dots, x^{n-t_8-1}A_8 \in \text{Span}(B)$. Now we show that $x^{t_5-t_7}A_7 \in \text{Span}(B)$. From Statement 2 of Lemma 2, we have

$$(3.10) \quad x^{t_5-t_7}A_7 = c_5vA_5 - q'_7(x)A_7 - q'_8(x)A_8,$$

where $\deg(q'_7(x)) < (t_5 - t_7)$. In the above discussion, we have shown that $x^iA_8 \in \text{Span}(B)$ for $1 \leq i \leq n - t_8 - 1$. Clearly, $q'_8(x)A_8 \in \text{Span}(B)$. And, also the term $c_5vA_5, q'_7(x)A_7 \in \text{Span}(B)$ (since $\deg(q_7(x)) < (t_5 - t_7)$). Therefore,

$x^{t_5-t_7}A_7 \in \text{Span}(B)$. As above, after putting the value of $x^{t_5-t_7}A_7$ in the equation obtained by multiplying Equation (3.10) by $x, x^2, \dots, x^{t_1-t_5-1}$, successively, we can show that $x^{t_5-t_7+1}A_7, x^{t_5-t_7+2}A_7, \dots, x^{t_1-t_7-1}A_7 \in \text{Span}(B)$. From Statement 4 of Lemma 2, we have

$$(3.11) \quad x^{t_1-t_7}A_7 = c_1vwA_1 - q'_7(x)A_7 - q'_8(x)A_8,$$

where $\deg(q'_7(x)) < t_1 - t_7$. As above, we can show that $x^{t_1-t_7}A_7, x^{t_1-t_7+1}A_7, \dots, x^{n-t_7-1}A_7 \in \text{Span}(B)$. Now we show that $x^{t_5-t_6}A_6 \in \text{Span}(B)$. From Statement 5 of Lemma 2, we have

$$(3.12) \quad x^{t_5-t_6}A_6 = c_5uA_5 - q'_6(x)A_6 - q'_7(x)A_7 - q'_8(x)A_8,$$

where $\deg(q'_6(x)) < (t_5 - t_6)$. In the above discussion, we have shown that $x^iA_j \in \text{Span}(B)$ for $1 \leq i \leq n - t_j - 1, j = 7, 8$. Clearly, $q'_7(x)A_7, q'_8(x)A_8 \in \text{Span}(B)$. Therefore, $x^{t_5-t_6}A_6 \in \text{Span}(B)$. In a similar way, as above, after putting the value of $x^{t_5-t_6}A_6$ in the equation obtained by multiplying Equation (3.12) by $x, x^2, x^3, \dots, x^{t_1-t_5-1}$, successively, we can show that $x^{t_5-t_6+1}A_6, x^{t_5-t_6+2}A_6, \dots, x^{t_1-t_6-1}A_6 \in \text{Span}(B)$. From Statement 7 of Lemma 2, we have

$$(3.13) \quad x^{t_1-t_6}A_6 = c_1uwA_1 - q'_6(x)A_6 - q'_7(x)A_7 - q'_8(x)A_8,$$

where $\deg(q'_6(x)) < t_1 - t_6$. Again as above, we can show that $x^{t_1-t_6+1}A_6, x^{t_1-t_6+2}A_6, \dots, x^{n-t_6-1}A_6 \in \text{Span}(B)$. Now we show that the next term $x^{t_1-t_5}A_5 \in \text{Span}(B)$. From Statement 8 of Lemma 2, we have

$$(3.14) \quad x^{t_1-t_5}A_5 = c_1wA_1 + q_5(x)A_5 + q_6(x)A_6 + q_7(x)A_7 + q_8(x)A_8,$$

where $\deg(q_5(x)) < t_1 - t_5$. In the above discussion, we have shown that $x^iA_j \in \text{Span}(B)$ for $1 \leq i \leq n - t_j - 1, j = 6, 7, 8$. Clearly, $q_6(x)A_6, q_7(x)A_7, q_8(x)A_8 \in \text{Span}(B)$. Therefore, $x^{t_1-t_5}A_5 \in \text{Span}(B)$ (since $\deg(q_5(x)) < t_1 - t_5$). Multiplying Equation (3.14) by $x, x^2, x^3, \dots, x^{n-t_1-1}$ and then putting the value of $x^{t_1-t_5}A_5$ in the equation obtained, we can show that the terms $x^{t_1-t_5+1}A_5, x^{t_1-t_5+2}A_5, \dots, x^{n-t_5-1}A_5 \in \text{Span}(B)$.

Case (2A). Let $t'_8 = t_4, t'_7 = t_3$ and $t'_6 = t_2$. Let $t'_4 = t_3$. As in Case 1, by using Statement 1 of Lemma 2 for $i = 4, 3$ and 1, successively, we can show that $x^{t_4-t_8}A_8, x^{t_4-t_8+1}A_8, \dots, x^{n-t_8-1}A_8 \in \text{Span}(B)$. Similarly, as in Case 1, by using Statements 3 and 4 of Lemma 2, successively, we can show that $x^{t_3-t_7}A_7, x^{t_3-t_7+1}A_7, \dots, x^{n-t_7-1}A_7 \in \text{Span}(B)$. Again, as in Case 1, by using Statement 6 and then Statement 7 of Lemma 2, successively, we can show that $x^{t_2-t_6}A_6, x^{t_2-t_6+1}A_6, \dots, x^{n-t_6-1}A_6 \in \text{Span}(B)$. In a similar fashion, as in Case 1, by using Statement 8 of Lemma 2, we can show that $x^{t_1-t_5}A_5, x^{t_1-t_5+1}A_5, \dots, x^{n-t_5-1}A_5 \in \text{Span}(B)$.

Case (2B). Let $t'_8 = t_4, t'_7 = t_3, t'_6 = t_2$ and $t'_4 = t_2$. As in Case 1, by using Statement 1 of Lemma 2 for $i = 4, 2$ and 1, successively, we can show that $x^{t_4-t_8}A_8, x^{t_4-t_8+1}A_8, \dots, x^{n-t_8-1}A_8 \in \text{Span}(B)$. Similarly, as in Case 1, by using Statements 3, 4, 6, 7 and then 8 of Lemma 2, successively, we can show that the rest of elements belongs to $\text{Span}(B)$.

Case (3). Let $t'_8 = t_6$, $t'_7 = t_3$ and $t'_6 = t_5$. As in Case 1, by using Statement 1 of Lemma 2 for $i = 6, 5$ and 1, successively, we can show that $x^{t_6-t_8}A_8, x^{t_6-t_8+1}A_8, \dots, x^{n-t_8-1}A_8 \in \text{Span}(B)$. Similarly, as in Case 1, by using Statements 3 and 4 of Lemma 2, successively, we can show that $x^{t_3-t_7}A_7, x^{t_3-t_7+1}A_7, \dots, x^{n-t_7-1}A_7 \in \text{Span}(B)$. Again, as in Case 1, by using Statement 5 and then Statement 7 of Lemma 2, successively, we can show that $x^{t_5-t_6}A_6, x^{t_5-t_6+1}A_6, \dots, x^{n-t_6-1}A_6 \in \text{Span}(B)$. In a similar fashion, as in Case 1, by using Statement 8 of Lemma 2, we can show that $x^{t_1-t_5}A_5, x^{t_1-t_5+1}A_5, \dots, x^{n-t_5-1}A_5 \in \text{Span}(B)$.

The remaining cases are as follows:

Case (4). If $t'_8 = t_7$, $t'_7 = t_3$ and $t'_6 = t_5$.

Case (5). If $t'_8 = t_7$, $t'_7 = t_5$ and $t'_6 = t_2$.

Case (6). If $t'_8 = t_7$, $t'_7 = t_5$ and $t'_6 = t_5$.

Case (7). If $t'_8 = t_6$, $t'_7 = t_3$ and $t'_6 = t_2$.

Case (8). If $t'_8 = t_6$, $t'_7 = t_5$ and $t'_6 = t_2$.

Case (9). If $t'_8 = t_6$, $t'_7 = t_5$ and $t'_6 = t_5$.

Case (10). If $t'_8 = t_4$, $t'_7 = t_5$ and $t'_6 = t_5$, **(10A):** $t'_4 = t_3$, **(10B):** $t'_4 = t_2$.

Case (11). If $t'_8 = t_4$, $t'_7 = t_3$ and $t'_6 = t_5$, **(11A):** $t'_4 = t_3$, **(11B):** $t'_4 = t_2$.

Case (12). If $t'_8 = t_4$, $t'_7 = t_5$ and $t'_6 = t_2$, **(12A):** $t'_4 = t_3$, **(12B):** $t'_4 = t_2$.

In a similar way as above, by using Lemma 2, we can show that B spans B_2 in these cases.

Now we show that B spans B_1 . From Equation (2.1), we have a homomorphism $\phi : C \rightarrow R_{u^2, v^2, p, n}$. Therefore, $C/\text{Ker}\phi \simeq \phi(C)$ and $\phi(C)$ is a cyclic code over the ring $R_{u^2, v^2, p}$. Thus, we have $C/\text{Ker}\phi$ as a cyclic code over $R_{u^2, v^2, p}$. Therefore, from Theorem 4.1 of [8], the minimal spanning set B_ϕ of the code $C/\text{Ker}\phi$ is $\{A_1 + \text{Ker}\phi, xA_1 + \text{Ker}\phi, \dots, x^{n-t_1-1}A_1 + \text{Ker}\phi, A_2 + \text{Ker}\phi, xA_2 + \text{Ker}\phi, \dots, x^{t_1-t_2-1}A_2 + \text{Ker}\phi, A_3 + \text{Ker}\phi, xA_3 + \text{Ker}\phi, \dots, x^{t_1-t_3-1}A_3 + \text{Ker}\phi, A_4 + \text{Ker}\phi, xA_4 + \text{Ker}\phi, \dots, x^{t_4-t_4-1}A_4 + \text{Ker}\phi\}$. To show B spans B_1 , we only show that $x^{t_1-t_2}A_2 \in \text{Span}(B)$. In a similar way, we can show that $x^{t_1-t_2+1}A_2, \dots, x^{n-t_2-1}A_2, \dots, x^{t_4-t_4}A_4, \dots, x^{n-t_4-1}A_4 \in \text{Span}(B)$. Since B_ϕ spans $C/\text{Ker}\phi$, we can write $x^{t_1-t_2}A_2 + \text{Ker}\phi$ as a $R_{u^2, v^2, p}$ linear combination of the elements of B_ϕ , i.e., $x^{t_1-t_2}A_2 + \text{Ker}\phi = \sum_{i=0}^{n-t_1-1} \alpha_{i1}(x^i A_1 + \text{Ker}\phi) + \dots + \sum_{i=0}^{t_4-t_4-1} \alpha_{i4}(x^i A_4 + \text{Ker}\phi)$, where $\alpha_{ij} \in R_{u^2, v^2, p}$. Thus, $x^{t_1-t_2}A_2 - (\sum_{i=0}^{n-t_1-1} \alpha_{i1}(x^i A_1) + \dots + \sum_{i=0}^{t_4-t_4-1} \alpha_{i4}(x^i A_4)) \in \text{Ker}\phi$. Since $\text{Ker}\phi = \text{Span}(B_2)$ and B spans B_2 , we get $x^{t_1-t_2}A_2 \in \text{Span}(B)$. Similarly, we can show that $x^{t_1-t_2+1}A_2, \dots, x^{n-t_2-1}A_2, \dots, x^{t_4-t_4}A_4, \dots, x^{n-t_4-1}A_4 \in \text{Span}(B)$. This shows that B spans B_1 . It is easy to see that any elements of the spanning set B can not be written as the linear combination of its preceding elements and other elements in the spanning set B . Here we only show that $x^{t_1-t_3-1}A_2$ can not be written as linear combinations of others element of spanning set B . The proof is similar for the rest. Suppose, if possible $x^{t_1-t_2-1}A_2$ can be written as

linear combinations of the others element of the spanning set B . Then we have

$$\begin{aligned} x^{t_1-t_3-1}A_3 &= \sum_{i=0}^{n-t_1-1} \alpha_{1i}x^iA_1 + \sum_{i=0}^{t_1-t_2-1} \alpha_{2i}x^iA_2 + \sum_{i=0}^{t_1-t_3-2} \alpha_{3i}x^iA_3 \\ &+ \sum_{i=0}^{t'_4-t_4-1} \alpha_{4i}x^iA_4 + \sum_{i=0}^{t_1-t_5-1} \alpha_{5i}x^iA_5 + \sum_{i=0}^{t'_6-t_6-1} \alpha_{6i}x^iA_6 \\ &+ \sum_{i=0}^{t'_7-t_7-1} \alpha_{7i}x^iA_7 + \sum_{i=0}^{t'_8-t_8-1} \alpha_{8i}x^iA_8, \end{aligned}$$

where $\alpha_{ji} = \beta_{j1}^{(i)} + u\beta_{j2}^{(i)} + v\beta_{j3}^{(i)} + uv\beta_{j4}^{(i)} + w\beta_{j5}^{(i)} + uw\beta_{j6}^{(i)} + vw\beta_{j7}^{(i)} + uvw\beta_{j8}^{(i)} \in \mathbb{F}_p$ (Note that i is not a power of β it is a notation). We have

$$\begin{aligned} &x^{t_1-t_3-1}(vf_3(x) + uvf_{3,4}(x) + wf_{3,5}(x) + uwf_{3,6}(x) + vwf_{3,7}(x) + uvwf_{3,8}(x)) \\ &= f_1(x) \sum_{i=0}^{n-t_1-1} \beta_{11}^{(i)}x^i + uf_1(x) \sum_{i=0}^{n-t_1-1} \beta_{12}^{(i)}x^i + uf_{1,2}(x) \sum_{i=0}^{n-t_1-1} \beta_{11}^{(i)}x^i \\ &+ uf_2(x) \sum_{i=0}^{t_1-t_2-2} \beta_{21}^{(i)}x^i + vf_1(x) \sum_{i=0}^{n-t_1-1} \beta_{13}^{(i)}x^i + vf_{1,3}(x) \sum_{i=0}^{n-t_1-1} \beta_{11}^{(i)}x^i \\ &+ vf_{2,3}(x) \sum_{i=0}^{t_1-t_2-1} \beta_{21}^{(i)}x^i + vf_3(x) \sum_{i=0}^{t_1-t_3-2} \beta_{31}^{(i)}x^i \\ &+ uvm_2(x) + wm_3(x) + uw(m_4(x) + vwm_5(x) + uvwm_6(x)), \end{aligned}$$

where $m_2(x), \dots, m_6(x)$ are polynomials in $\mathbb{F}_p[x]$. By comparing both sides, we have $\beta_{11}^{(i)} = 0$, $\beta_{12}^{(i)} = 0$ for $0 \leq i \leq n - t_1 - 1$, $\beta_{21}^{(i)} = 0$ for $0 \leq i \leq t_1 - t_2 - 1$, and $x^{t_1-t_3-1}f_3(x) = f_1(x) \sum_{i=0}^{n-t_1-1} \beta_{13}^{(i)}x^i + f_3(x) \sum_{i=0}^{t_1-t_3-2} \beta_{31}^{(i)}x^i$. Note that $\deg(x^{t_1-t_3-1}f_3(x)) = t_1 - 1$ but $\deg(f_1(x) \sum_{i=0}^{n-t_1-1} \beta_{13}^{(i)}x^i) \geq t_1$ and $\deg(f_3(x) \sum_{i=0}^{t_1-t_3-2} \beta_{31}^{(i)}x^i) \leq t_1 - 2$. Hence, this gives a contradiction. \square

Theorem 7. *Let n be a positive integer relatively prime to p and C be a cyclic code of length n over the ring $R_{u^2, v^2, w^2, p}$. If $C = \langle f_1(x) + uf_2(x), vf_3(x) + uvf_4(x), w(f_5(x) + uf_6(x)), w(vf_7(x) + uvf_8(x)) \rangle$ with $t_i = \deg(f_i(x))$, $1 \leq i \leq 8$, and $t'_7 = \min\{t_3, t_5\}$, then C has rank $n + t_1 + t'_7 - t_3 - t_5 - t_7$. The minimal spanning set B of the code C is $B = \{f_1(x) + uf_2(x), x(f_1(x) + uf_2(x)), \dots, x^{n-t_1-1}(f_1(x) + uf_2(x)), vf_3(x) + uvf_4(x), x(vf_3(x) + uvf_4(x)), \dots, x^{t_1-t_3-1}(vf_3(x) + uvf_4(x)), w(f_5(x) + uf_6(x)), xw(f_5(x) + uf_6(x)), \dots, x^{t_1-t_5-1}w(f_5(x) + uf_6(x)), w(vf_7(x) + uvf_8(x)), xw(vf_7(x) + uvf_8(x)), \dots, x^{t'_7-t_7-1}w(vf_7(x) + uvf_8(x))\}$.*

Proof. The proof is similar to the above theorem. \square

4. Minimum distance

Let n be a positive integer not relatively prime to p . Let C be a cyclic code of length n over $R_{u^2, v^2, w^2, p}$. From Eq. (2.9), we have $C_8 = \{f(x) \in \mathbb{F}_p[x] \mid uvwf(x) \in C\} = \langle f_8(x) \rangle$. Also, we know that C_8 is a cyclic code over \mathbb{F}_p .

Theorem 8. *Let n be a positive integer not relatively prime to p . If $C = \langle A_1, A_2, \dots, A_8 \rangle$ is a cyclic code of length n over $R_{u^2, v^2, w^2, p}$. Then $w_H(C) = w_H(C_8)$.*

Proof. Let $M(x) = m_0(x) + um_1(x) + vm_2(x) + uvm_3(x) + wm_4(x) + uvm_5(x) + vvm_6(x) + uvwm_7(x) \in C$, where $m_0(x), m_1(x), \dots, m_7(x) \in \mathbb{F}_p[x]$. We have $uvwM(x) = uvwm_0(x)$, $w_H(uvwM(x)) \leq w_H(M(x))$ and $uvwC$ is subcode of C with $w_H(uvwC) \leq w_H(C)$. Thus $w_H(uvwC) = w_H(C)$. Therefore, it is sufficient to focus on the subcode $uvwC$ in order to prove the theorem. Since $w_H(C_8) = w_H(uvwC)$, we get $w_H(C) = w_H(C_8)$. \square

Definition. Let $m = b_{l-1}p^{l-1} + b_{l-2}p^{l-2} + \dots + b_1p + b_0$, $b_i \in \mathbb{F}_p$, $0 \leq i \leq l-1$, be the p -adic expansion of m .

- (1) If $b_{l-i} \neq 0$ for all $1 \leq i \leq q$, $q < l$, and $b_{l-i} = 0$ for all $i, q+1 \leq i \leq l$, then m is said to have a p -adic length q zero expansion.
- (2) If $b_{l-i} \neq 0$ for all $1 \leq i \leq q$, $q < l$, $b_{l-q-1} = 0$ and $b_{l-i} \neq 0$ for some $i, q+2 \leq i \leq l$, then m is said to have p -adic length q non-zero expansion.
- (3) If $b_{l-i} \neq 0$ for $1 \leq i \leq l$, then m is said to have a p -adic length l expansion or p -adic full expansion.

Lemma 3. *Let C be a cyclic code over $R_{u^2, v^2, w^2, p}$ of length p^l where l is a positive integer. Let $C = \langle f(x) \rangle$ where $f(x) = (x^{p^{l-1}} - 1)^b h(x)$, $1 \leq b < p$. If $h(x)$ generates a cyclic code of length p^{l-1} and Hamming distance d , then the Hamming distance $d(C)$ of C is $(b+1)d$.*

Proof. For $c \in C$, we have $c = (x^{p^{l-1}} - 1)^b h(x)m(x)$ for some $m(x) \in \frac{R_{u^2, v^2, w^2, p}[x]}{(x^{p^l} - 1)}$. Since $h(x)$ generates a cyclic code of length p^{l-1} , we have

$$\begin{aligned} w(c) &= w((x^{p^{l-1}} - 1)^b h(x)m(x)) \\ &= w(x^{p^{l-1}b} h(x)m(x)) + w({}^b C_1 x^{p^{l-1}(b-1)} h(x)m(x)) + \dots \\ &\quad + w({}^b C_{b-1} x^{p^{l-1}} h(x)m(x)) + w(h(x)m(x)). \end{aligned}$$

Thus, $d(C) = (b+1)d$. \square

Theorem 9. *Let C be a cyclic code over $R_{u^2, v^2, w^2, p}$ of length p^l , where l is a positive integer. Then, $C = \langle A_1, A_2, \dots, A_8 \rangle$ where $f_1(x) = (x-1)^{t_1}$, $f_2(x) = (x-1)^{t_2}$, \dots , $f_8(x) = (x-1)^{t_8}$ for some $t_1 > t_2, t_3 > t_4 > t_8 > 0$, $t_2 > t_6$, $t_3 > t_7$ and $t_1 > t_5 > t_6, t_7 > t_8 > 0$.*

- (1) If $t_8 \leq p^{l-1}$, then $d(C) = 2$.
- (2) If $t_8 > p^{l-1}$, let $t_8 = b_{l-1}p^{l-1} + b_{l-2}p^{l-2} + \cdots + b_1p + b_0$ be the p -adic expansion of t_8 and $f_8(x) = (x-1)^{t_8} = (x^{p^{l-1}} - 1)^{b_{l-1}}(x^{p^{l-2}} - 1)^{b_{l-2}} \cdots (x^{p^1} - 1)^{b_1}(x^{p^0} - 1)^{b_0}$.
 - (a) If t_8 has a p -adic length q zero expansion or full expansion ($l = q$), then $d(C) = (b_{l-1} + 1)(b_{l-2} + 1) \cdots (b_{l-q} + 1)$.
 - (b) If t_8 has a p -adic length q non-zero expansion, then $d(C) = 2(b_{l-1} + 1)(b_{l-2} + 1) \cdots (b_{l-q} + 1)$.

Proof. The first claim easily follows from Theorem 2. From Theorem 8, we see that $d(C) = d(C_8) = d(\langle (x-1)^{t_8} \rangle)$. Hence, we only need to determine the minimum weight of $C_8 = \langle (x-1)^{t_8} \rangle$.

(1) If $t_8 \leq p^{l-1}$, then $(x-1)^{t_8}(x-1)^{p^{l-1}-t_8} = (x-1)^{p^{l-1}} = (x^{p^{l-1}} - 1) \in C$. Thus, $d(C) = 2$.

(2) Let $t_8 > p^{l-1}$. (a) If t_8 has a p -adic length q zero expansion, we have $t_8 = b_{l-1}p^{l-1} + b_{l-2}p^{l-2} + \cdots + b_{l-q}p^{l-q}$, and $f_8(x) = (x-1)^{t_8} = (x^{p^{l-1}} - 1)^{b_{l-1}}(x^{p^{l-2}} - 1)^{b_{l-2}} \cdots (x^{p^{l-q}} - 1)^{b_{l-q}}$. Let $h(x) = (x^{p^{l-q}} - 1)^{b_{l-q}}$. Then $h(x)$ generates a cyclic code of length p^{l-q+1} and minimum distance $(b_{l-q} + 1)$. By Lemma 3, the subcode generated by $(x^{p^{l-q+1}} - 1)^{b_{l-q+1}}h(x)$ has minimum distance $(b_{l-q+1} + 1)(b_{l-q} + 1)$. By induction on q , we can see that the code generated by $f_8(x)$ has minimum distance $(b_{l-1} + 1)(b_{l-2} + 1) \cdots (b_{l-q} + 1)$. Thus, $d(C) = (b_{l-1} + 1)(b_{l-2} + 1) \cdots (b_{l-q} + 1)$.

(b) If t_8 has a p -adic length q non-zero expansion, we have

$$t_8 = b_{l-1}p^{l-1} + b_{l-2}p^{l-2} + \cdots + b_1p + b_0, b_{l-q-1} = 0.$$

Let $r = b_{l-q-2}p^{l-q-2} + b_{l-q-3}p^{l-q-3} + \cdots + b_1p + b_0$ and $h(x) = (x-1)^r = (x^{p^{l-q-2}} - 1)^{b_{l-q-2}}(x^{p^{l-q-3}} - 1)^{b_{l-q-3}} \cdots (x^{p^1} - 1)^{b_1}(x^{p^0} - 1)^{b_0}$. Since $r < p^{l-q-1}$, we have $p^{l-q-1} = r + j$ for some non-zero j . Thus, $(x-1)^{p^{l-q-1}-j}h(x) = (x^{p^{l-q-1}} - 1) \in C$. Hence, the subcode generated by $h(x)$ has minimum distance 2. By Lemma 3, the subcode generated by $(x^{p^{l-q}} - 1)^{b_{l-q}}h(x)$ has minimum distance $2(b_{l-q} + 1)$. By induction on q , we can see that the code generated by $f_8(x)$ has minimum distance $2(b_{l-1} + 1)(b_{l-2} + 1) \cdots (b_{l-q} + 1)$. Thus, $d(C) = 2(b_{l-1} + 1)(b_{l-2} + 1) \cdots (b_{l-q} + 1)$. \square

Let C be a cyclic code of length n over the ring $R_{u^2, v^2, w^2, p}$. Let d_L be the Lee distance of a code C . By Theorem 1 of [9], we have

$$\left\lceil \frac{d_L - 1}{8} \right\rceil \leq n - \log_p s |C|.$$

If C is a free code, then we get

$$\left\lceil \frac{d_L - 1}{8} \right\rceil \leq n - \text{Rank}(C).$$

5. The Gray map

Let w_L and w_H denote the Lee weight and the Hamming weight respectively. We define the Lee weight as follows:

$$w_L(\alpha) = w_H(\phi_L(\alpha)) \text{ for all } \alpha \in R_{u^2, v^2, w^2, p},$$

where the Gray map $\phi_L : R_{u^2, v^2, w^2, p} \rightarrow \mathbb{F}_p^8$ is defined as follows:

$$\begin{aligned} & \phi_L(\alpha_1 + u\alpha_2 + v\alpha_3 + uv\alpha_4 + w\alpha_5 + uw\alpha_6 + vw\alpha_7 + uvw\alpha_8) \\ &= (\alpha_8, \alpha_6 + \alpha_8, \alpha_7 + \alpha_8, \alpha_4 + \alpha_8, \alpha_5 + \alpha_6 + \alpha_7 + \alpha_8, \alpha_2 + \alpha_4 + \alpha_6 + \alpha_8, \\ & \quad \alpha_3 + \alpha_4 + \alpha_7 + \alpha_8, \alpha_1 + \alpha_2 + \alpha_3 + \alpha_4 + \alpha_5 + \alpha_6 + \alpha_7 + \alpha_8). \end{aligned}$$

The Gray map naturally extend to $R_{u^2, v^2, w^2, p}^n$ as distance preserving isometry $\phi_L : (R_{u^2, v^2, w^2, p}^n, \text{ Lee weight}) \rightarrow (\mathbb{F}_p^{8n}, \text{ Hamming weight})$ as follows

$$\phi_L(\alpha_1, \alpha_2, \dots, \alpha_n) \rightarrow (\phi_L(\alpha_1), \phi_L(\alpha_2), \dots, \phi_L(\alpha_n)), \quad \forall \alpha_i \in R_{u^2, v^2, w^2, p}.$$

By linearity of the map ϕ_L we obtain the following theorem.

Theorem 10. *If C is a linear code over $R_{u^2, v^2, w^2, p}$ of length n , size p^k and minimum lee weight d , then $\phi_L(C)$ is a p -ary linear code with parameters $[8n, k, d]$.*

5.1. Gray images of cyclic codes over $R_{u^2, v^2, w^2, p}$

Definition. Let T be the cyclic shift such that $T(c_0, c_1, \dots, c_{n-1}) = (c_{n-1}, c_0, \dots, c_{n-2})$. Then a linear code C is an l -quasicyclic code of length n if it is invariant under T^l .

Lemma 4. *Let T be the cyclic shift. Then*

$$\phi \circ T = T^8 \circ \phi_L.$$

Proof. Let $\tilde{r} = (r_0, r_1, \dots, r_{n-1}) \in R_{u^2, v^2, w^2, p}^n$ and ϕ_L be the Gray map defined in 2.2. Then

$$(\phi \circ T)(\tilde{r}) = (\phi_L(r_{n-1}), \phi_L(r_0), \dots, \phi_L(r_{n-2})).$$

We also know that

$$\phi_L(r_0, r_1, \dots, r_{n-1}) = (\phi_L(r_0), \phi_L(r_1), \dots, \phi_L(r_{n-1})),$$

where each $\phi_L(r_i)$ is of length 8. Therefore, if we apply the cyclic shift eight times, the whole of $\phi_L(r_{n-1})$ will shift from the end to the beginning, which means we will get

$$(T^8 \circ \phi_L)(\tilde{r}) = (\phi_L(r_{n-1}), \phi_L(r_0), \dots, \phi_L(r_{n-2})).$$

Hence, we get the result. \square

Theorem 11. *Let C be a cyclic code of length n over the ring $R_{u^2, v^2, w^2, p}$. Then $\phi_L(C)$ is a 8-quasicyclic binary linear code of length $8n$ over \mathbb{F}_p .*

Proof. Let C be a cyclic code over $R_{u^2, v^2, w^2, p}$. Then we know that $T(C) = C$. Now applying ϕ_L to both sides we get

$$\phi_L(T(C)) = \phi_L(C).$$

But, by Lemma 4, we know that $\phi \circ T = T^8 \circ \phi_L$. Therefore,

$$\phi_L(C) = \phi_L(T(C)) = (\phi_L \circ T)(C) = T^8(\phi_L(C))$$

which implies that $\phi_L(C)$ is invariant under T^8 . Hence, $\phi_L(C)$ is a 8-quasicyclic code. \square

6. Examples

Example 1. Cyclic codes of length 4 over the ring $R_{u^2, v^2, w^2, 2}$. We have

$$x^4 - 1 = (x - 1)^4 \text{ over } \mathbb{F}_2.$$

Let $g = x - 1$, some of the non zero cyclic codes of length 4 over the ring $R_{u^2, v^2, w^2, 2}$ with generator polynomials, rank and minimum distance are given in Table 1.

TABLE 1. Non zero cyclic codes of length 4 over $R_{u^2, v^2, w^2, 2}$.

Non-zero generator polynomials	Rank	$d(C)$
$\langle vwg^2 + (c_0 + c_1x)uvw \rangle$	2	2
$\langle vwg + c_0uvw \rangle$	3	2
$\langle uwg^3 + c_1vvg^3 + c_0uvw^2 \rangle$	1	4
$\langle uwg^3 + c_0uvw(c_2 + c_3x), vvg^3 + (c_0 + c_1x)uvw \rangle$	3	4
$\langle wg^3 + c_2uwg^2 + c_3vvg^2 + c_4uvw, uvw^2 \rangle$	2	2
$\langle wg + c_1uw + c_2vw \rangle$	3	2
$\langle uwg^3 + c_1uwg^3 + c_2vvg + c_3uvw, vvg^2 + c_0uvw \rangle$	3	2
$\langle vg^3 + c_1uv + c_1vvg, uwg^2 + c_0wg, wg^2, uv, vw \rangle$	8	1
$\langle g^2 + v + u + c_1w, vg + u, ug + c_1w, uv + c_1w, wg, uv, vw \rangle$	8	1

Example 2. Cyclic codes of length 3 over the ring $R_{u^2, v^2, w^2, 3}$. We have

$$x^3 - 1 = (x - 1)^3 \text{ over } \mathbb{F}_3.$$

Let $g = x - 1$, some of the non zero cyclic codes of length 3 over the ring $R_{u^2, v^2, w^2, 3}$ with generator polynomials, rank and minimum distance are given in Table below:

Example 3. Cyclic codes of length 5 over the ring $R_{u^2, v^2, w^2, 5}$. We have

$$x^5 - 1 = (x - 1)^5 \text{ over } \mathbb{F}_5.$$

Let $g = x - 1$, some of the non zero cyclic codes of length 5 over the ring $R_{u^2, v^2, w^2, 5}$ with generator polynomials, rank and minimum distance are given in Table below:

TABLE 2. Non zero cyclic codes of length 3 over $R_{u^2, v^2, w^2, 3}$.

Non-zero generator polynomials	Rank	$d(C)$
$\langle uwg^2 + vwg, vwg^2 \rangle$	3	3
$\langle uwg + c_2vw + c_1uvw, vwg + c_0uvw \rangle$	4	2
$\langle wg + c_1uw + c_2vw, uwg + c_0vw \rangle$	3	2
$\langle vwg + c_2uw + c_1vw, wg + c_0uw \rangle$	5	1
$\langle vg^2 + c_4uwg + c_3w + c_2uw, uwg + wg^2 + c_1uw, uwg^2 + c_0vw, vwg, uvw \rangle$	5	1
$\langle ug + c_3v + c_2wg + c_1vw, vg + c_1wg + c_1vw, uv + c_0wg + c_1vw + c_0vw, wg^2 + c_0vw + c_0uvw \rangle$	8	3
$\langle g^2 + c_3u + c_0uvw, ug + c_2w + c_0uvw, v + c_1w + c_0uvw, wg + c_0uvw, uwg + c_0uvw \rangle$	6	1

TABLE 3. Non zero cyclic codes of length 5 over $R_{u^2, v^2, w^2, 5}$.

Non-zero generator polynomials	Rank	$d(C)$
$\langle vwg^4 + uwg^3 \rangle$	1	5
$\langle uwg^4 + vwg^3 + uvwg, vwg^4 + uvwg^2, uvwg^3 \rangle$	3	4
$\langle uwg^3 + c_4vwg + uvw(c_2 + c_3x), vwg^2 + (c_0 + c_1x)uvw \rangle$	5	3
$\langle wg^4 + uw(c_6 + c_7x) + vwc_8 + uvwc_9, uwg^2 + vw(c_4 + c_5x) + uvw(c_2 + c_3x), vwg^2 + uvw(c_0 + c_1x) \rangle$	5	3
$\langle uvw^4 + wc_2g^3 + ww(c'_3 + c'_4x) + vw(c'_0 + c'_1x) + uvw, wg^4 + c_9uwg + vw(c_6 + c_7x) + c_8uvw, uwg^2 + vw(c_3 + c_4x + c_5x^2) + uvw(c_1 + c_2x), vwg^2 + uvw(c_0 + c_1x) \rangle$	6	3
$\langle uvw^3 + c'_4wg^2 + c'_3uw + c'_2vw, wg^4 + uwg + c'_1vw + c'_0uvw, uwg^3 + vw(c_3 + c_4x) + c_5uvw, vwg^3 + (c_0 + c_1x)uvw, uvwg^2 \rangle$	6	3
$\langle vg^4 + c'_0uwg + w(c_1 + c'_2x) + c'_3uw + c'_4vw, uwg^2 + c'_7wg^2 + c'_8uw + c'_9vw + c'_5uvw, wg^4 + c_4uwg^2 + c_5vw + c_6uvw, uwg^3 + vw + uvwg, vwg + (c_0 + c_1x)uvw \rangle$	8	2

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